

Extending The Network Lifetime Using Optimized Energy Efficient Cross Layer Module (OEEXLM) In Wireless Sensor Networks

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Abstract

In wireless sensor network, the primary design is to save the energy consumption as much as possible while achieving the given task. Most of recent researches works have only focused on the individual layer issues and ignore the importance of inter working between different layers in a sensor network. In this paper, we use a cross-layer approach to propose an energy-efficient and extending the life time of the sensor network. This protocol which uses routing in the network layer, and the data scheduling in MAC layer. The main objective of this paper is to provide a possible and flexible approach to solve the conflicts between the requirements of large scale, long life-time, and multi-purpose wireless sensor networks. This OEEXLM module gives better performance compared to all other existing protocols. The performance of OEEXLM module compared with S-MAC and directed diffusion protocol.

Index terms: Routing, Medium access control, life time of the network, energy efficiency, OEEXLM module, Wireless Sensor Networks.

1. Introduction

Wireless sensor network consists of large number of sensor nodes are randomly distributed in a region. Each node has a limited energy supply and generates information when ever event occurs that needs to be communicated to a sink node. The focus of this paper is on the computation of optimal energy usage for data transfer and link schedule that maximize the network lifetime. However most wireless sensor network energy consumption operations involve sensing, computing, and communicating [1]. Generally, communication between nodes consumes more energy than local processing or collecting data operation. The geographical nature of the deployment space of nodes makes quasi impossible the replacement or the recharging operations of batteries. The challenge is to economize energy inside every node in order to maintain as long as possible the network functionality.

2. Related Works

Many research works are developed for energy efficiency at each layer of protocol stack by proposing

new algorithms and protocols. In particular, MAC layer was of great interest for many researchers because it is considered as an important source of energy wastage such as overhearing, collision, control packet overheads and idle listening [2].

In order to decrease or if possible to eliminate these various sources of energy wastage, several protocols has been proposed during last years and which are divided into two main classes: Schedule based protocol TDMA and Contention based protocol.

In TDMA protocols [3] are employed to avoid collisions by associating a slot time for each sensor node in a given cluster. This protocols more complex in the WSN where the nodes in general have a same priority and very limited resources. The Contention-based protocols are known as CSMA-based are usually used in the multi-hop wireless networking. But this protocol generates collisions due to these useless retransmissions which cause energy consumption wastage and time consuming in data transmission.

Some of the power control problems are discussed in [4,5]. Approaches at MAC layer are Dermikol et al [4] infers that some pros and cons of the some of the existing protocol. One of the first attempts at MAC protocol for WSN is PAMAS, which reserves battery power intelligently powering off users that are not actively transmitting or receiving packets. This protocol decreased the energy consumption of the network but maximum latency occurs. Raghavendra et al [5] the Power Aware Medium Access protocol and Signaling (PAMAS) is a CSMA based protocol in which the nodes that are not actively transmitting or receiving should power themselves off. The protocol requires the nodes to have two separate channels (control and data), which will require two radios at each node increasing the cost, size and complexity of the sensor design.

Existing MAC protocol [6,7] turn off the transceiver when there is no communication between the nodes also the power savings in these papers results reducing the idle listening power but also decreasing the collisions. In these papers they proposed adaptive listening incurs overhearing. Sleep and listen periods are predefined and constant, which decreases the efficiency of the algorithm under variable traffic load.

A new generation of MAC protocols that is Cross-layer MAC protocols using several layers in order to

optimize energy consumption has been emerged. These layers can be divided into interaction mode or unification mode. In the interaction mode, the MAC protocol is built by exploiting the data generated by other adjacent layers.

Approaches at network layer in WSNs are mostly used to implement the routing of the incoming data, as quoted in [8]. It is known that generally in multihop networks the source nodes cannot reach the sink directly. Therefore, intermediate sensor nodes have to relay their packets. Generally, the implementation of routing tables offers the solution. These routing tables contain lists of node options for any given packet destination Definition of the routing tables is the task of the routing algorithm along with the help of the routing protocol for their construction and maintenance.

MAC-CROSS Protocol [9] is an example of Cross-layer approach which allows the interaction between MAC and information of the network layers by making only the communicating nodes in listening mode and by putting other nodes into sleep mode. In order to avoid collisions, MAC-CROSS uses the control messages RTS/CTS/ACK. On the other hand, a Cross-layer design mode by unification requires the development of only one layer including at the same time functionalities of considered layers.

In this paper, a unified cross layer module XLM[10] is developed which achieves efficient and reliable event communication between the nodes with minimum energy expenditure. But in this protocol the end-to-end delay increases for low value duty cycle. When the duty cycle is low $\delta = 0.1$, 14% of the transmitted packets are dropped due to retransmission timeout. Because sender nodes cannot find any neighbors that satisfy the constraints.

Weiyang Ge et al[11] addressed the rate optimization for multicast communications at the media access control (MAC) layer, and explore transport layer erasure coding to enhance multicast reliability in wireless sensor networks. In this approach they are not addressed the energy consumption and end-to-end delay.

3. Proposed OEEEXLM Module

The proposed cross layer approach replaces the entire traditional layer protocol architecture that has been so far used in WSNs. In this protocol we integrate the medium access control and routing to improve the performance of the network. The communication in OEEEXLM module is based on initiative concept. This module allows the each node to decide whether it can participate in communication or not. Consequently, this protocol uses adaptive receiving, stagger scheduling algorithm and logical link decision algorithm.

A node starts a transmission by transmitting to its neighborhood an RTS packet to indicate that it has a packet to send. Upon receiving an RTS packet, each

neighborhood node i decide to participate to communication or not, that decision can be determined by an initiative “ I ” defined as follows:

$$I = \begin{cases} \beta & \text{if } \left. \begin{array}{l} TA > R + T \\ \gamma = \gamma_{th} \\ E_{rev} = E_{rev}^{thd} \end{array} \right\} \\ 0 & \text{otherwise} \end{cases} \quad (1)$$

The initiative is set to 1 if all four conditions are true or satisfied. Out of four conditions, the first two conditions related to MAC layer and next two conditions related to network layer. The first condition checks event β occurs or not. If $\beta = 1$ then the remaining three conditions checked by the node, otherwise the nodes are go to sleep mode.

Because of this, idle listening and overhearing are avoided. The second condition is adaptive receiving scheme, the receiver node wait for TA seconds if the node does not receive any packets from upper layer it goes to sleep state. The third condition is the node will check the buffer size γ_{th} if the size is less than threshold it will choose the alternate path to reach the destination. Fourth condition is that source node from upper layer checks the lower layer receiver node energy if it is greater than threshold E_{rev}^{thd} then it will transmit packet otherwise it will choose the alternate path.

Using this initiative concept, OEEEXLM module overcomes collision, overhearing, idle listening and improve the throughput, link reliability and extend the life time of the network.

3.1 Basic terms Used in OEEEXLM

The following assumption can be considered for OEEEXLM protocol to study the performance. The network topology used in this protocol is grid architecture. Data flows from n layer to $n-1$ layer until the data packet reached the destination. For simulation MICA2 mote specification is considered. In this protocol the value of duty cycle is denoted by δ and is defined as ratio of the time a node is active. The duty cycle is varied with respect to data transmission. The sleep time for each node is T_{sleep} sec. The listen period of each node is less than or equal to transmission period of upper layer. Transmission period each layer is 60ms. With in 60ms if lower layer node does not receive any data from upper layer it will wait for TA sec and go to sleep.

3.2 Initiation of Transmission

When the event occurs that is $\beta = 1$ layer n node has a data packet to transmit, it sends RTS signal to the lower layer, the lower layer sends the CTS signal to the upper layer with response the CTS signal the layer n sends

the data to the layer ($n-1$) for 60ms. During the receiving period, if a node senses that the channel is idle and its neighbours are not communicating for a time TA , it will go to sleep until its subsequent sending periods. After receiving the packet from n layer, the $n-1$ layer route the packets to the destination, each source node determines a path to the destination by selecting a lower layer node under its coverage randomly. Thus messages flow in the correct direction, but do not use the same path every time. Thus this data exchange scheme provides both collision avoidance and reliable retransmission.

4. Mechanism of OEEEXLM

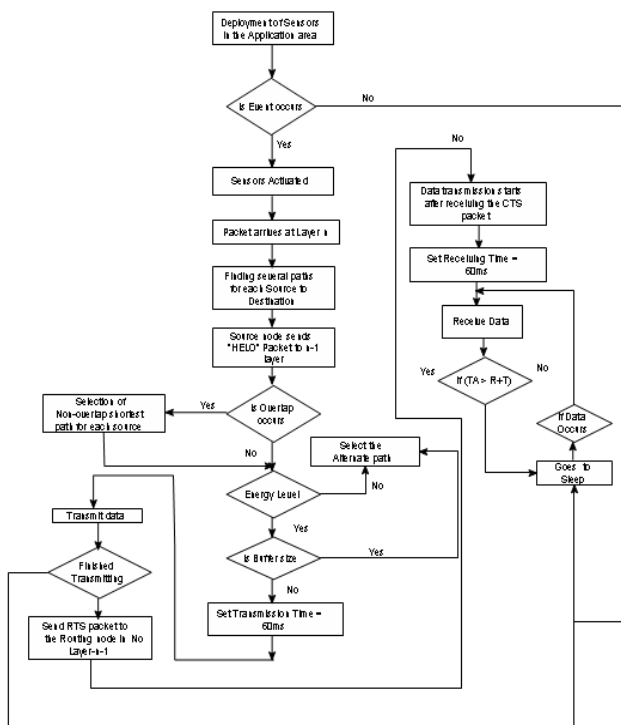


Figure1 Mechanism of OEEEXLM

Figure1 explains the mechanism of OEEEXLM module. In this module we integrate the MAC and network layer. In MAC layer we proposed Power Efficient MAC protocol which is used to overcome idle listening, collisions, Hidden terminal problem, and also to provide a low latency compared to other existing MAC protocols. The PE-MAC uses the three algorithms to achieve the energy efficiency.

4.1 Clock Synchronization algorithm

In clock synchronization algorithm layer n comprises the source nodes. Initially the nodes in layer n are in sending mode for 60ms while those in the layer ($n-1$) are in receiving mode. The remaining nodes are in

sleep state. Assume that the source nodes generate packets (layer- n).

As the next layer nodes are in the receiving mode, the source nodes can transmit the packets directly without checking the status of the lower layer nodes. The packets are stored in the buffer space of the lower layer nodes. Once the receiving period of layer ($n-1$) nodes ends, the sending period for layer ($n-1$) starts. The layer ($n-2$) nodes shift from sleep mode to receive mode. The remaining nodes enter the sleep mode.

4.2 Logical Link Decision Algorithm

The LLD algorithm is implemented in this protocol to ensure that two source nodes do not transmit the packets to the same receiver at the same time. This algorithm is implemented initially when each source node determines a link to the sink. When all the source nodes determine their links, the links are compared to ensure that there is no overlap in the existing links. If there is an overlap, the LLD algorithm uses a logical link decision to get the new link to the sink.

4.3 Adaptive Receiving

In Adaptive Receiving algorithm, layer n sends RTS packet to the lower layer $n-1$, the lower layer sends the CTS signal to the upper layer n . With response the CTS signal from layer $n-1$, the layer n sends the data to the layer $n-1$ for 60ms. Adaptive receiving scheme employs a time interval TA to handle traffic varies.

During the receiving period, if a node in the $n-1$ layer senses that the channel is idle and data from n layer are not communicating for a time TA , it will go to sleep. Suppose within 60ms if the layer $n-1$ senses any data from the upper layer n it wakes up and receive the data. This adaptive receiving scheme reduces the packet drop and improves the throughput of the network.

In network layer we proposed two algorithms: a congestion control and alternate path algorithm. These two algorithms are used to improve the lifetime of multihop sensor network by avoiding the collision.

4.4 Alternate path algorithm

In alternate path algorithm, each has routing table. Before route the data packet each source finds a multiple path to the sink. In this each source node initiate HELLO message to all lower layer nodes.

The HELLO message contains source ID, type of node whether it is sink or intermediate node and energy level which is shown in figure 2 .

Source ID	Type of Node	Energy level
2 Byte	1 Byte	4 Byte

Figure 2 HELLO Message Packet Format

Each node in the lower layer updates its table once it receives the *HELLO* message. If any node in lower layer gets more than one *HELLO* message it sends the negative acknowledgement *NACK* signal to the corresponding source node. Then that particular source node chooses the alternate path to reach the sink.

After setting an alternate route the nodes in the routing path check its energy level of receiving node. If the energy level is less than the threshold E_{rev}^{thd} again the node chooses the alternate path by sending *HELLO* message and the routing table is updated.

4.5 Congestion Control Mechanism and Algorithm

The congestion in the network is due to two traffic. One is due to generated packet that is whenever node detects the events it will generate the data packets and that is to be transmitted to the destination through intermediate nodes. The rate of generated packets at the node i is denoted by λ_i .

Since the network is a multi hop each node plays a dual role that is it acts as source as well as router. During the transmission each node in lower layer receives the data packet from the upper layer until the data packet reaches the destination. These packets are referred to as relay packets. If the $n-1$ layer node receives the data packets from layer n the rate of relay packet of node i of $n-1$ layer is $\lambda_{i,n}^n @1$. The input rate of the buffer for the node i depends on the rate of relay packets $\lambda_{i,n}^n @1$ and the rate of generated packets λ_i . In the OEEXLM module the rate of input packet at node i 's buffer α_i , can be represented as

$$\alpha_i = \lambda_i + \lambda_{i,n}^n @1 \tag{2}$$

The node is active for a fraction of duty cycle $\delta = 60ms$. Hence the average time taken to transmit and receive a data packet can be given by

$$T_{tx} = 1 + e_i \alpha_i A_{pkt} \tag{3}$$

$$T_{rx} = \lambda_{i,n}^n @1 t_{pkt} \tag{4}$$

Where t_{pkt} is the average time taken to successfully transmit a packet to another node and e_i is the error packet rate.

The proposed OEEXLM module avoids packet drops due to congestion by not allowing upper layer nodes to transmit data packets if there is not available buffer size γ . This can be controlled by a congestion control algorithm. A lower layer node changes the path of transmission based on its buffer status. During transmission, the lower layer node i allows P_i packets to be transmitted by the upper layer nodes. Figure 3 shows the buffer queue model of the node.

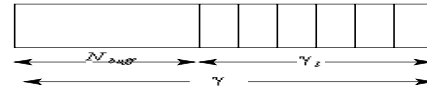


Figure 3 Buffer Queue Model

Given the condition that the lower layer node has a proportionally higher probability to access the medium, even then the lower layer node may not be able to forward all P_i packets. So, in the next transmission, the lower layer node will have some packets from the previous transmission, which might cause congestion within a few successive transmissions. Therefore, for each transmission upper layer nodes check the routing table if the buffer size γ is less than the buffer threshold γ_{thd} the upper layer node takes the alternate path. The maximum number of packets held by the buffer is N_{buff} . Transmission only occurs when $\gamma > \gamma_{thd}$

$$\gamma_{thd} = N_{buff} \geq \sum_{i=0}^{i=n} P_i \tag{5}$$

Where $N_{buff} = \gamma @ \gamma_l$

As a result of the congestion control algorithm, the OEEXLM module avoids layer-to-layer congestion in the networks. Because of the energy consumed by the node is reduced there by improving the life time of the network.

The medium control access scheme uses three algorithms to improve energy efficiency and in the network layer we use two algorithms to improve the life time of the networks. These two layers are integrated by using initiative conditions as per equation (1).

5. Network Topology

Figure 4 shows the data flow for the OEEXLM module. The nodes in the network operate in three different modes- sleeping, receiving and sending. Each node goes to sleep periodically to save energy and then wakes up and listens to see if any other node wants to talk to it. During sleep, the node turns off its radio; therefore the energy waste due to idle listening can be reduced.

The nodes with the same layer-count are given the same schedule, and the sending and receiving periods are staggered layer by layer such that when one node is in sending mode, its lower-layer node is in receiving mode. After receiving a message from an upper-layer node, each node can transmit it to the lower-layer node in the subsequent sending period. So a packet can be transmitted to sink nodes through multi-layers fleetly and the end-to-end latency is reduced.

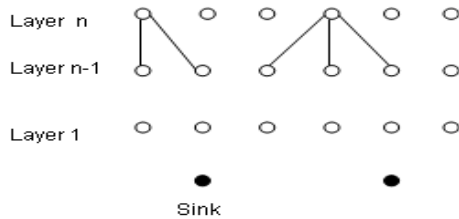


Figure 4 Data flow Diagram for OEEXML module

5.1 Performance Evaluation

The Table 1 shows the simulation setup used for simulation.

Parameters	Value
Transmission Range	250 m
Network Area	100 x 100
Number of Sensors	100- 1500
Packet rate	5 pkt/sec
Packet size	512 bytes
Radio Bandwidth	76kbps
Transmitting Power	75mW (270J)
Receiving Power	36mW (129.6J)
Power Consumption in Sleep mode	100μ W (0.36 J)
Sending and Receiving Slot	50msec
Type of mote	Mica2
Initial energy of sensor node	2KJ
Energy Threshold E^{thd}	0.001mJ

Table 1 Simulation Setup

5.1.1 Latency Vs Network area (MAC)

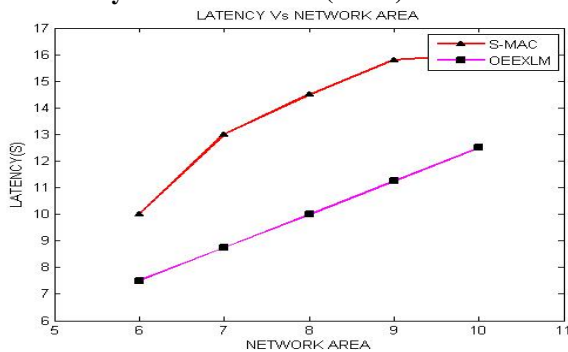


Figure 5 Latency Vs Network Area

Figure 5 infers that in S-MAC a node has to wait till its neighbouring node as to awake to transmit the message to it. This results there will be some amount of delay in S-MAC protocol, which is absent in OEEXML module. In OEEXML module whenever a node is in send state, the lower layer node is in receiving state. So the source node can transmit the message to the lower layer without checking whether the lower layer node is listening or

sleeping mode. Hence the source node can transmit its message fleetly to the sink through multi layers.

5.1.2 Energy Consumption Vs Network Area (MAC)

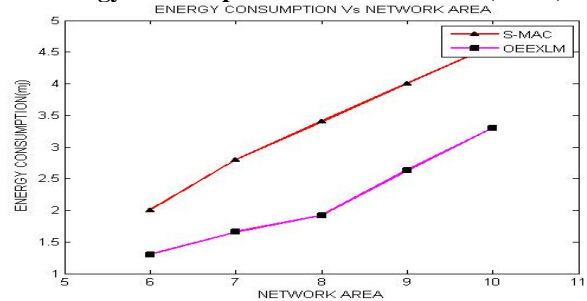


Figure 6 Energy Consumption Vs Network Area

The figure 6 infers that depicts the energy consumption of S-MAC and OEEXML module. The comparison is made for a simulation setup with ten layers, ten nodes in each layer and a four sink. The graph shows that OEEXML module protocol uses less energy than S-MAC. This is because the idle listening protocol dominates the energy consumption in S-MAC but TA can make OEEXML module go to sleeping mode earlier, and the energy consumption is reduced. The energy consumption of OEEXML module is 35% -83% les than existing S-MAC protocol.

5.1.3 Latency Vs Number of Hop (MAC)

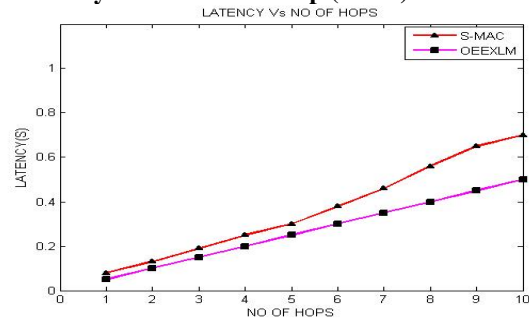


Figure 7 Latency Vs Number of Hops

The Figure 7 shows that the latency encountered in OEEXML module compared with that in S-MAC with respect to number of hops. Obviously the latency of OEEXML module is lesser than that in the S-MAC because of staggered scheduling algorithm. In S-MAC a node has to wait till its neighbour is awake to transmit the message to it. This problem is not occurring in OEEXML module. In OEEXML module whenever a node is in send state, the lower layer node is in receiving state. So the source node can transmit the message to the lower layer without checking whether the lower layer node is listening. Hence the source node can transmit its message fleetly to the sink through multi layers.

5.1.4 Latency Vs Number of Hops (Routing)

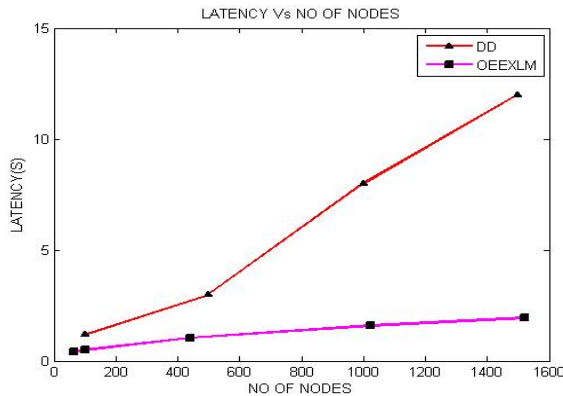


Figure 8 Comparison of Latency Vs Number of hops with DD

Figure 8 infers the delay as a function of the number of nodes in the WSN. The delay increases with the number of hops increasing. Our simulation shows the average delay of the proposed protocol is better than that of the directed diffusion. This is because in OEEEXLM module, for each transmission each node in the routing checks the buffer condition, so that no queue of data occurs in the buffer. Due to alternate path algorithm the shortest delay occur compared to other schemes. As we expected, data packets are routed through different node-with the help of proper design of algorithm in routing. Hence, the network congestion can be avoided.

5.1.5 Energy consumption Vs Number of Nodes (Routing)

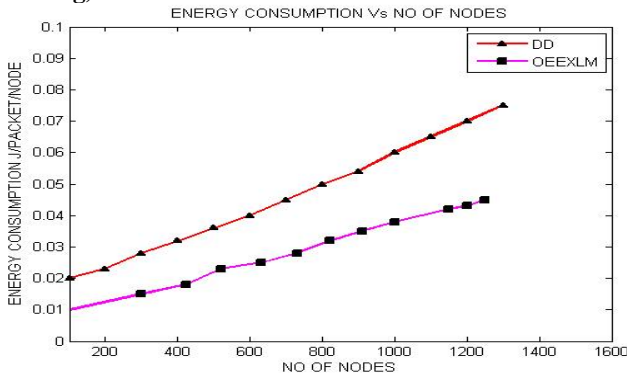


Figure 9 Energy consumption Vs Number of nodes

Figure 9 infers that energy consumption per node versus number of nodes. The value of node energy consumption gives the average energy dissipated by the node in order to transmit the packet from source to drain. The same metric is used in [6] to determine the energy efficiency level of WSNs. It is calculated as follows:

$$Node\ Energy\ Consumption = \frac{\sum_{i=1}^N e_i,initial @ e_i,resi}{N \times \sum_{j=1}^S data P_j} \quad (6)$$

Where N is the total number of nodes, $e_i,initial$ initial energy of nodes, $e_i,resi$ is the residual energy of the nodes, S is the number of sinks and P_j is the number of data packets received by the sink j .

The simulation result infers that there is lower node energy consumption in CRSL protocol over the other schemes. The energy consumption of nodes in OEEEXLM module is 34% to 84% lesser than when compared with directed diffusion. This results shows that the energy efficiency of OEEEXLM module is stable and has little impact by the increase of the network size, while the performance of other schemes degrades with larger network size.

5.1.6 Network Life Time

The system lifetime is defined as the number of rounds for which 75% of the nodes are still alive. CRSL The transmitted and received energy costs for the transmission of a k-bit data message between two nodes separated by a distance of r meters are given by Eqs.(7)and(8), respectively.

$$E_t k,r = E_{tx} k + E_{amp} r^a k \quad (7)$$

$$E_r k = E_{rx} k \quad (8)$$

In Eqn (7) $E_t k,r$ denotes the total energy dissipated in the transmitter of the source node, while $E_r k$ in Eqn. (8) represents the energy cost incurred in the receiver of the destination node. Parameters E_{tx} and E_{rx} are per bit energy dissipation for transmission and reception, respectively, and $E_{amp} r^a$ denotes the energy required by the transmitted amplifier to maintain an acceptable radio for transferring data reliably. The free-space propagation model is applied, and the transmit amplifier $E_{amp} r^a$ is given by Eqn.8

$$E_{amp} = \epsilon_{FS} r^2 \quad (9)$$

where ϵ_{FS} is the transmitted amplifier parameter. The set of parameter given in [7,8]. $\epsilon_{FS} = 10 \text{ pJ/b/m}^2$.

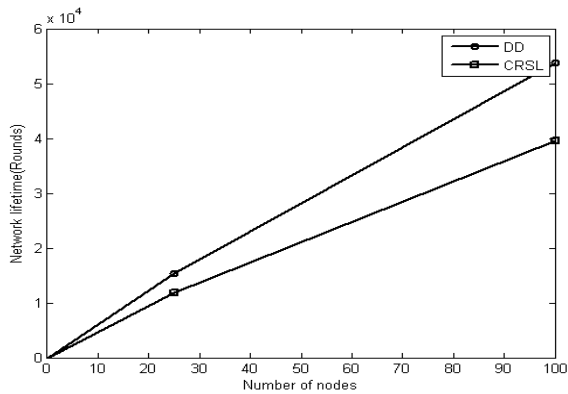


Figure 10 Network lifetime

From Figure 10, we infer that the network lifetime is increased by using the alternate path and congestion control algorithm. The network lifetime is given in terms of rounds till which 75% of nodes are alive. A network is assumed to be useless when one of the sensor’s energy is below the threshold.

5.1.7 Data Packet delivery ratio

Data packet delivery ratio can be calculated as the ratio between the number of data packets that are sent by the source and the number of data packets that are received by the sink. It can be denoted as R

$$Data\ Delivery\ Ratio = \frac{Sum\ of\ data\ packets\ received\ at\ sink}{Required\ data}$$

This parameter R indicates both efficient of the routing protocol and the effort required to receive data. In the ideal scenario the ratio should be equal to 1. If the ratio falls significantly below the ideal ratio, then it could be an indication of some of the packet dropped because of faults in the protocol design. However, if the ratio is higher than the ideal ratio, then it is an indication that the sink receives a data packet more than once. It is not desirable because reception of duplicate packets consumes the more energy. The relative number of duplicates received by the sink is also important because based on that number the sink, can possibly take an appropriate action to reduce the redundancy.

Figure11 shows the data packet delivery ratio of DD and OEEXLM modules. To eliminate packet loss we use a rate of 5 packets / second. It is found that the delivery ratio of the two protocols increase as the node density increases. When node density is high, there are more nodes available for data forwarding, and this increases the delivery ratio. Directed diffusion protocol offers less packet delivery rates, compared to OEEXLM module because it does not adapt well its behavior to network size increase. The OEEXLM module has maintained constant delivery rates throughout the

simulated scenarios because the paths are selected based on the energy availability and buffer size

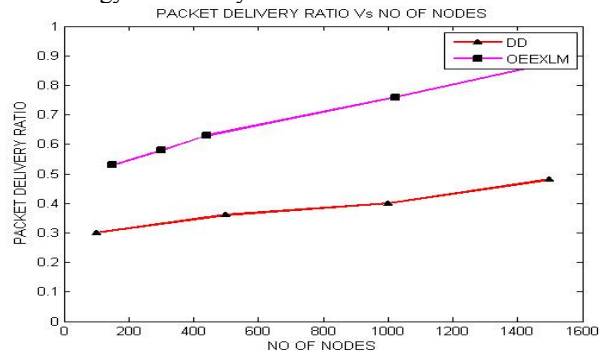


Figure 11 Packet delivery Ratio Vs Number of Nodes

5.2 Performance Variations with respect to Buffer size

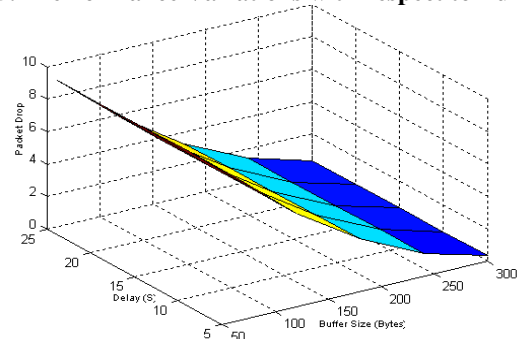


Figure12 Performance Variations with respect to Buffer size

Figure12 shows the performance variations when the buffer size varies from 50 bytes to 300 bytes. When the buffer size is 50 bytes number of packets reached at the destination is minimum because of congestion but packet delay is minimum. When the buffer size 300 bytes the number of packets dropped is minimum due to large buffer and the average delay per packet increases due to the increased queuing delay.

5.2.1 Packet loss with respect Buffer Size

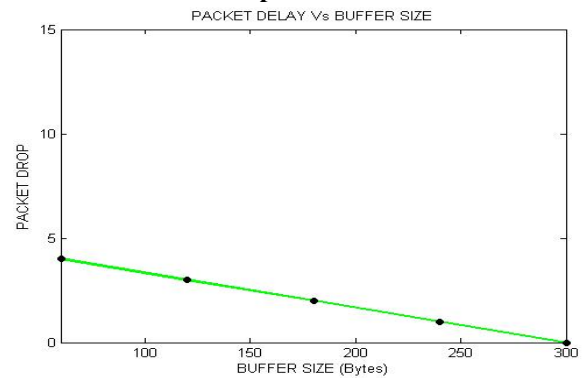


Figure 13 Packet loss Vs Buffer size

Figure 13 infers that packet loss get reduced due to increasing buffer size. The number of packets dropped due to buffer overflow in the case of the OEEXML module almost zero. This is because each time after receiving the packets from upper layer the buffer size is calculated and updated in the routing table. Depending on updated value the routing path algorithm send the HELLO message to the other nodes in the routing.

5.2.2 Packet Delay with Buffer size

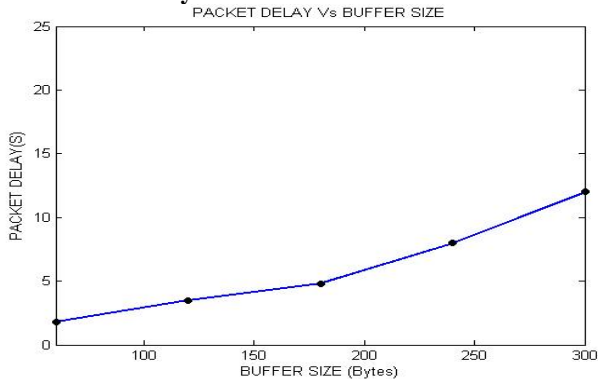


Figure 14 Packet Delay Vs Buffer Size

Figure 14 shows Packet delay Vs buffer size. The graph infers that as the buffer size increases the average delay per packet increases due to the increased queuing delay. When the buffer size is 120 bytes, packet delay is only 4 second from source to destination. When the buffer size is 300 bytes length the delay is 12 seconds because of queuing delay. From this graph, we find that as the buffer size increases the packet delay also increases. But throughput is maximum because the packet drop is minimum by compromising the delay.

6. Conclusion and Future work

In this paper we presented the Cross-Layer Design to improve the performance of the wireless sensor networks. This protocol design is used to give the direct interactions between the Network layer and the MAC layer. The traditional Network layer and MAC layer have been removed, thus simplifying the protocol stack. Simulation results that our proposed scheme has higher node energy efficiency, lower average delay and control overhead than the directed diffusion protocol and S-MAC protocol. The energy consumption of nodes in OEEXML module is 34% to 84% lesser than when compared with directed diffusion. Further the network life time is 78% improved compared to DD protocol. In future we are going to extend the OEEXML design to Physical layer.

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