Highly Efficient Kerberos Style Authentication and Authorization for Cloud Computing

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Summary

This article presents kerberos system which achieves security goals for the cloud users like authentication, integrity, confidentiality and it provides distributed authentication services across the insecure network. This kerberos style authentication is based on the secret key technology which keeps online password attackers away and it is implemented by using Windows Active Server.

General Terms:

Mathematics Subject Classification (2000), 68-04

Key words:

Security, Kerberos, Cloud Computing, Authentication, Authorization

1. Introduction

Cloud computing is an emerging, on-demand and internet based technology used by global customers to improve their business performance. Cloud is highly scalable, flexible and platform independable which has security loopholes like attacks, data loss, other authentication and security issues. Bo Wang [1] stated that to utilize cloud services by authorized user and to secure cloud data, it is necessary to use secure authentication system. This technology consists of dedicated user stations and distributed servers which provide services to multiple users and require the ability to accurately identify the cloud user. The authentication service can be achieved by using Kerberos which requires cloud users to prove identity for each service invoked, the servers prove their identity to cloud users. Kerberos provide reliable authentication over open and insecure networks where communication within the hosts belonging to it may be intercepted.

Gagan Dua[2] presents an improved method which prevents replay attacks and password attacks by using Triple Password Scheme. Three passwords are stored on Authentication Server. Authentication server sends two passwords to Ticket Granting Server (one of the Application Servers) by encrypting with the secret key shared between Authentication Server and Ticket Granting Server. Similarly Ticket Granting Server sends one password to application server by encrypting with the secret key shared between TGS and application server. Meanwhile service granting ticket is transferred to users by encrypting it with the password that TGS receives from AS which helps to prevent replay attack.

Trapti Ozha[8] designed authentication protocol based on secret key encryption technology, uses data encryption algorithm for encryption to extend the security for the system. Ved.M.Kshirsagar [9] proposed a scheme which achieves batch auditing where multiple delegated auditing different users can be performed from tasks simultaneously by the third party auditor- a privacy preserving manner. Pathan Mohd Shafi [3] designed a system which combines text based password authentication with cued click point and Kerberos authentication protocol. In this system, password consists of sequence of images in which user selects one click point per image and selects a sound signature. The sound signature is validated for accessing the system. Sulochana.V[5] proposed the sequence level to access the cloud services. The sequence level authentication generates password at five levels and then concatenates into one single password.

Shubha Bharill[4] designed an authentication model for cloud based on Kerberos protocol using threshold cryptography to provide more security and to increase the availability of key. This model filters unauthorized access and reduce the burden of computation and memory usage of cloud provider against authentication checks for each client. Sulochana.V[6] also presented location based authentication in which cloud user chooses a location as password and its longitude and latitude are extracted by GPS device, it is validated by local server and the cloud user gets authenticated and start accessing the cloud services. Sulochana.V[7] developed a puzzle based authentication scheme in which cloud user registers and solves the puzzle, puzzle solving time and sequence of image block are stored and validated by local server and the cloud user gets authenticated and start accessing the cloud services. The proposed Kerberos stvle authentication works on the basis of ticket to allow nodes communicating over a non secure network to prove their identity to one another in a secure manner.

Manuscript received November 5, 2016 Manuscript revised November 20, 2016

2. Kerberos Style of Authentication and Authorization

2.1 Mathematical Model

Let A1 be the set of cloud users in the system

A1 = {CU, Auth_Srvr, Ti_Gr_Srvr, CS} CU→Cloud User, Auth_Srvr→Authentication Server, Ti_Gr_Srvr→ Ticket Granting Server CS→ Cloud Server on which user wants to access the cloud services

Let A2 be the set of objects in the system A2={ID_{CU}, ID_{Ti_Gr_Srvr}, T_S, K_{CU}, K_{CS}, Lifetime, AD_{CU}, Ticket_{Ti_Gr_Srvr}, Ticket_{CS}} ID_{CU}→ID of Cloud User, ID_{Ti_Gr_Srvr}→ID of Ticket Granting Server K_{CU}→Key of Cloud User K_{CS}→Key of Cloud Server AD_{CU}→ Network Address of Cloud User

$$\begin{split} & \text{Ticket}_{Ti_Gr_Srvr} = E(K_{Ti_Gr_Srvr} ~ [~ K_{Ti_Gr_Srvr} ~ \| ~ ID_{CU} ~ \| ~ AD_{CU} ~ \| \\ & \text{ID}_{Ti_Gr_Srvr} ~ \| ~ TS_2 ~ \| ~ \text{lifetime]}) \\ & \text{Ticket}_{CS} = E(~ K_{CS} ~ [~ K_{CU,CS} ~ \| ~ ID_{CU} ~ \| ~ AD_{CU} ~ \| ~ ID_{CS} ~ \| ~ TS_4 ~ \| \\ & \text{lifetime]}) \\ & \text{Authenticator} = E(K_{CU,Ti_Gr_Srvr} ~ [ID_{CU} ~ \| ~ AD_{CU} ~ \| ~ TS]) \end{split}$$

Three Stages of Problem Solving

a) Authentication service exchange is to obtain Ticket Granting Ticket (Ticket_{Ti_Gr_Srvr}) **Step1:** CU \rightarrow Auth_Srvr ID_{CU}||ID_{Ti_Gr_Srvr}||TS₁ **Step2:** Auth_Srvr \rightarrow CU E(K_{CU}[K_{CU,Ti_Gr_Srvr} || ID_{Ti_Gr_Srvr} || TS₂ || lifetime || Ticket_{Ti_Gr_Srvr}]) Cloud User requests Auth_Srvr for ticket to access Ti_Gr_Srvr. Auth_Srvr checks Cloud user authenticity and sends Ticket_{Ti_Gr_Srvr} to cloud user.

b) Ticket granting service exchange is to obtain server granting ticket (Tickets $_{cs}$)

Step3: CU→Ti_Gr_Srvr

 $ID_{cs} \parallel Ticket_{Ti_Gr_Srvr} \parallel Authenticator_{CU}$

Step4: $Ti_Gr_Srvr \rightarrow CU$ $E(Kcu_{Ti_Gr_Srvr} [K_{CU,CS} || ID_{CS} || TS_4 || Ticket_{CS}])$

Cloud user sends the ID and $Ticket_{Ti}Gr_Srvr}$ to Ti_Gr_Srvr to authenticate itself. Ti_Gr_Srvr replies with a $Ticket_{cs}(Ticket to Cloud Server)$

c) Cloud user|Cloud server authentication exchange is to obtain service

Step5: CU \rightarrow CSTicket_{CS} || AuthenticatorStep6:CS \rightarrow CUE (K_{CU,CS} [TS₅₊₁]) for mutualauthentication. Cloud User sends Ticket_{CS} to Cloud serverto authenticate and access service.

2.2 Conventions

In this paper new terms are used and elaborated as follows,

Key Distribution Center: KDC is the heart of Kerberos realm, handles the distribution of keys and tickets. It provides Kerberos authentication services by issuing encrypted tickets which require secret keys to decode.

Session Key: Temporary private keys generated Kerberos which is known to cloud user to encrypt the communication between cloud user and cloud server.

Authentication Server (AS): An authentication service is to know the password of cloud user and stores in centralized database, shares a unique secret key with TGS and Cloud Server.

Ticket Granting Server (TGS): This server issues service tickets to cloud user upon request.

Ticket Granting Ticket (TGT): The ticket is issued by Authentication Server, TGT is encrypted in cloud user password which is known only by the user and KDC.

2.3 Working Model

This kerberos style provides strong authentication for client/server applications by using secret key cryptography and verify the identities of network server on open network without relying on authentication by host operating system, physical security of the host on the network.When cloud user wants to access a service from cloud server which requires Kerberos ticket, the server grants access subscribed service to cloud user based on the ticket. In the first phase, cloud user logs on to a workstation, cloud process running in the workstation issues a message to AS. Authentication Server validates username, password and issues TGT, session key to the cloud user which can be used to communicate with the server. TGT and session key are encrypted using key generated from the cloud user password. Cloud user decrypts the message and gets TGT used for communication with the TGS.

In the second phase, cloud user presents TGT to the TGS along with an authenticator. The authenticator includes ID of the cloud user and timestamp. TGT is encrypted with a secret key that is shared by AS and TGS. TGS can decrypt the message and gets the session key which decrypts the authenticator that is received from the cloud user. TGS checks the sender details by validating TGT with the details in the Authenticator and incoming packet network address. If details are validated, TGS issues a Service Granting Ticket to the cloud user and encrypted using the secret key shared between TGS and the Cloud Server. TGS generates a session key shared between the cloud user and cloud server for secure communication, sends the session key to the cloud server by encrypting with the ticket. The TGS sends session key to cloud user by encrypting the message with the session key shared between the cloud user and TGS. In the third phase, cloud user presents ticket to the cloud server along with the authenticator. The cloud user encrypts the message by using session key that is sent by TGS. Cloud server uses session key to encrypt the message from the cloud user. If credentials of the cloud user are correct, the cloud server will issue a response to the cloud user in case of a mutual authentication is required.

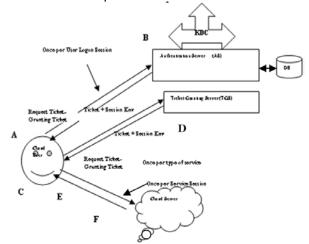


Fig 1. Working Model of the Kerberos Style Authentication

A→ Cloud User logs on to workstation and request service on host B→ AS verifies user's access right in database, creates ticket-granting ticket and session key. Results are encrypted using key derived from cloud user's password C→ Workstation prompts user for password and uses password to decrypt incoming message, then sends ticket and authenticator that contains cloud user's name, network address and time to TGS D→TGS decrypts tickets and authenticator verifies request then creates tickets for requested server E→ Workstation sends ticket and authenticator to server F→ Cloud server verifies that ticket and authenticator match, then grants access to service. If mutual authentication is required, cloud server returns authenticator

The Kerberos style authentication and authorization can be implemented by Windows Active Server. Kerberos system is based on the distribution of tickets. A ticket is an encrypted data generated by the Kerberos server and it is used by a service. It gives the right for the cloud user to access the service, assuming the service can decrypt the ticket. It contains ticket's expiry date field. KINIT command line utility on Linux Operating System is used to obtain granting ticket from Kerberos server. This Kerberos style authentication provides three level of security. The first level is provided by session key which is renewed every time by the cloud user who requires to access new service stored on the cloud server. The second level of security is provided by the cloud user password which is stored in the database at TGS. The third level of security is provided by the cloud user password which has already been forwarded to the cloud server.

Screenshot for getting Tickets:

| [ediqa@psrlxediqa1 cshrc]\$ kinit -c /tmp/del/s.tmp sulokerbuser@INFAKRB.INFAQA.COM | | |
|---|-------------------|--|
| Password for sulokerbuser∉INFAKRB.INFAQA.COM: | | |
| [ediqa@psrlxediqa1 cshrc]\$ klist -c /tmp/del/s.tmp | | |
| Ticket cache: FILE:/tmp/del/s.tmp | | |
| Default principal: sulokerbuser@INFAKRB.INFAQA.COM | | |
| Valid starting | Expires | Service principal |
| 07/15/14 16:41:55 | 07/16/14 02:41:58 | krbtgt/INFAKRB.INFAQA.COM@INFAKRB.INFAQA.COM |
| renew until 07/16/14 16:41:55 | | |
| [ediqa@psrlxediqa1 cshrc]\$ | | |

3. Conclusion

Kerberos is a powerful authentication system that is transparent to the cloud user except when entering initial password. The Kerberos system provides authentication and strong cryptography to secure information systems across an entire network or enterprise and a highly effective solution to network security problems.

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