

A Novel Hybrid Algorithm Based on Word and Method Ranking for Password Security

Berker Tasoluk^{1†} and Zuhail Tanrikulu^{2††},

Istanbul University

Bogazici University

Abstract

It is a common practice to use a password in order to restrict access to information, or in a general sense, to assets. Right selection of the password is necessary for protecting the assets more effectively. Password finding/cracking try outs are performed for deciding which level of protection do used or prospective passwords offer, and password cracking algorithms are generated. These algorithms are becoming more intelligent and succeed in finding more number of passwords in less tries and in a shorter duration. In this study, the performances of possible password finding algorithms are measured, and a hybrid algorithm based on the performances of different password cracking algorithms is generated, and it is demonstrated that the performance of the hybrid algorithm is superior to the base algorithms.

Keywords:

Password cracking, password security, hybrid, brute force, dictionary attacks, password testing, information security, it security.

1. Introduction

It is a common practice to use a password in order to restrict access to information, or in a general sense, to assets. Right selection of the password is necessary for protecting the assets more effectively. Password finding/cracking try outs are performed for deciding which level of protection do used or prospective passwords offer, and password cracking algorithms are generated. These algorithms are becoming more intelligent and succeed in finding a greater number of passwords in less tries and in a shorter duration.

In this study, the performances of possible password finding algorithms are measured, and a hybrid algorithm based on different algorithms is generated. The study comprises of following sections: In the related work section, general concepts about information security are examined; and the studies found on literature about password finding/cracking are briefly summarized.

In the methodology section, the hybrid algorithm generated is described.

In the results section, the performances of different password cracking algorithms are measured, a hybrid algorithm based on these algorithms is generated, and it is demonstrated that the performance of the hybrid algorithm is superior to the base algorithms.

In the conclusion and future work section, the results are summarized and suggestion for possible future research is specified.

2. Related Work

2.1. The 'Information Security' Concept

When information security is discussed, three features of information security show up. The first is confidentiality, which is the information will only be accessed by authorized ones. Integrity is another feature, and can be regarded as, information will only be changed by authorized ones. The last one is availability, can be summarized as being able to access information when needed. Those three features are regarded as the 'CIA Triad'. Those three features can contradict with each other, enhancing one feature may adversely impact other information security features [1, 2]

Data is very valuable asset for companies. Losing confidentiality of data may cost companies 100s of millions of dollars or more to companies [3]. According to IBM 2020 Cost of Data Breach Report, average data loss costs companies \$3.86 Million [4]. There are numerous cases in which cyber attackers have successfully infiltrated data from companies [5-7].

We can recall several incidents relating to sensitive data breaches. Adobe had an incident in 2013 and lost between 38-153 million accounts [5, 8]. In 2012 incident, LinkedIn is estimated to lose over 100 Million accounts [9]. In another incident, Marriott Hotels lost 500 Million customer data [10]. Other examples include Uber breach in 2016 which affected

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57 million users [11] and Friendfinder breach in 2017 which affected 412 million user accounts [12].

An incident happened in Epsilon is estimated to cost \$4Billion [13]. Another good example regarding data security is cracking the Enigma machine code which is used by Germans in World War II [14]. This progress shifted the course of the war. Nowadays, military and national security systems are electronic based, the responsible parties must take all necessary precautions to protect these systems. Some protection mechanisms are added to systems in order to protect information. When a person or a program requests the data, the first step is, they must tell who they are (identification). Identification is an assertion of the requestor, the system must use extra mechanisms to validate if this assertion is correct. Those controls comprise of 'authentication' step. After authentication step, the system determines what can the requestor do on the system and data (read/write/delete/edit/etc.), this step is called 'authorization' [15, 16].

In the authentication step, three different 'factors' can be used. Those factors are identified as; something you know, something you have and something you are. If two of those options are used while authenticating users, it is said 'two factor authentication' is used. Likewise, when all three of those options are used, 'three factor authentication' is used. Both two and three-factor authentication are called 'strong authentication' schemes [16].

An example of 'something you know' which is used on authenticating users is passwords. Apart from this; birth place, birthdate, birth place, identification number, home address, cellular phone number, etc. are all examples of this authentication factor. Tokens, cellular phones, credit/debit cards are all examples of 'something you have' factor. The password or push notification sent to cellular phone, the digital signature embedded in cellular sim card, or the dynamic number found on the soft token application are generally used in authentication as a second factor, on top of 'something you know' [17]. Fingerprint, retina scan, iris scan, palm scan, voice recognition, signature are all examples of the third factor, 'something you are'.

2.2. Passwords

Passwords are used extensively on user authentication. The passwords chosen by users may not always provide adequate protection. Research shows that a significant amount of user passwords can be found in short times [18]. A research made by Bonneau which examines 70 million Yahoo users reveals that user passwords provide 20-bit security on an offline attack [19].

The process of finding an unknown password is called 'password cracking'. Password cracking is used by attackers in order to gain unauthorized access to systems, and also used by information security professionals to decide the security level of a system, or to gain legitimate access to a system whose password is lost / forgotten.

Password cracking can be online or offline. In online password cracking, different passwords are tried through connecting to remote server's service. This service can be anything which allows remote connection: FTP, HTTP, IMAP, rlogin, SSH, NTP, VNC, Remote Desktop, Samba, etc. are a few service / protocol examples which can be victim of password cracking attacks [20]. This attack can be detected and reacted by the system owners. The system may lock out the user after a specific number of failed logon attempts, or disable the system for a specified period, or may request additional information like 'Captcha' [21] or another factor of authentication (probably something you have, by sending an sms to cellular phone or sending push notification to an application on phone). In the offline attack, typically the hash of the password is gained, and password cracking is attempted to determine the password which gives the same hash value as the gained password hash.

2.3. Hash Function

The passwords are typically not stored as cleartext in systems databases, they are hashed, and hash values are stored. There are many different hash algorithms, some of the widespread ones being used are MD5 and SHA1 [22]. Hash functions provide a fixed output independent of the size of the input. Even one bit change in input may result significant changes in hash output. Hash functions are one-way, one can not decide input by just looking at the output. MD5 produces 128-bit output, while SHA-1 produces 160-

bit output. Below is a table which summarizes some hash algorithms and their respective outputs for the input word 'berker':

Example outputs of some hashing algorithms.

Algorithm / Hash	Hash
MD2	DB95F5FB244301298AF06A98EF73F0DA
MD4	A958EF83DB1CA9627C28A506B9B35988
MD5	B9B209C7129285852F61F807F332725A
SHA-1	F1FA12439950A7CB0FE46EE5B8911ADD C8E94104
SHA-2 (256)	FC6635A10054A2533935191F43A6906710 3FE33492288B454 98BC6E49312A9CD
SHA-2 (384)	D7952C2446A6893FE3AB25A68518A1000 65FAE0246CA305
SHA-2 (512)	077E91ABDD68291145FDD666569EFEBCEC1BE C36325FE7B57
RIPEMD-160	D74BA971799C91E23B78A9DE42E294CF 9C2B9AA7714FDA 1AD43DD0BD55A57525C3D9EB690B1653C0482 DB95E31434 FE4329C51F6B73CB3D36C3A1E43BB30120D
LM	870BE7B49860AF40
NT	F70AC408F8922D97038256067D18BABB
MySQL323	7BED7155063CEB94
MySQLSHA1	B1234211B767DE84C67277071FDCCEC0 B055DCF2
Cisco PIX	wHJNRpyB.b4CKMp5
VNC Hash	EEEE9CB65095D1C2
Base64	YmVya2Vy

2.4. Password Cracking Attacks

To crack a password, different approaches can be employed. These can be summarized as, brute force attacks, dictionary attacks, hybrid attacks and pre computed hash (rainbow) tables [23]. In the brute force attack, all possible combinations are tried sequentially. In the dictionary attack, dictionary files

are used, and the items in the dictionary are used. The dictionaries can be language words, names, most used passwords, etc. Hybrid attack is the combination of brute force and dictionary attack. A dictionary is used for base words, and some rules are applied to those words, and then those modified words are used in the attack. In the rainbow table attack, pre-computed hash tables are used, only a lookup to the is performed, if the hash value of the target password is on the table, then the password is cracked at that time, if not, then this attack is said to be not successful [24-26]. Some rainbow tables come free within a password cracking software, and some are sold separately. There are also some paid services that give you the original source value of the given hash.

Many methods are advised to increase the efficiency and effectiveness of password cracking attacks. Chou et al. used a dictionary based on character placement on keyboard, and cracked 114% more passwords compared to standard dictionary [27]. Schweitzer et al. used keyboard patterns and was able to find 2/11 more passwords than standard dictionary attacks [28]. Weir et al. modeled the used passwords using context-free grammars, created a dictionary and cracked 28% to 129% more passwords compared to John the Ripper program [29, 30]. Narayanan and Shmatikov used Markov models, which is used for language processing, for password cracking and got better results than Oechslin's rainbow attack [31]. Simon used Markov chains and cracked 80% of the passwords in less than 24 hours [32].

3. Methodology

3.1. Uniqueness of the study

The referenced studies propose one method to improve the performance of password cracking, and compare the performance with standard dictionary attacks or John the Ripper password cracking program [30]. In this research, one aim is to propose a method which compares the success criteria of the dictionary and word generation / modification rules. Another aim is to create a custom dictionary that will be used in the hybrid attack, by combining many different dictionaries and word generation / modification rules, whose success criteria is better than them all. The proposed success criteria and dictionary / rule combination method constitute the novelty brought to the field.

3.2. Structure of the hybrid algorithm

There are two success criterion for the password cracking algorithm. One is how early the algorithm managed to crack the password. The less tries the algorithm performs before revealing the password, the better its performance is. Let's name this number as r. If the algorithm can not crack the password, it will not try forever, after a certain number of times, it will try to crack the next password. Let's name maximum try number, as q. If we try to crack p different passwords, then our total try count becomes:

$$r(t) = r(1) + r(2) + \dots + r(p).$$

This is our first success criteria. The lesser this number is, more successful the algorithm is. Second criteria is what percentage of the passwords are cracked in first q try. Let's call this number as s:

$$s = \text{passwords revealed} / \text{total passwords}$$

This number is anything between 0 and 1, and can be shown as percentage as well. The greater this number is, the more successful the algorithm is. After having those numbers for different algorithms, we have two performance numbers for every one of them. We will first normalize those numbers among the algorithms, let's call the normalized versions as r(t,n) and s(n). Then we will weigh those numbers according to needs of the password cracking operation. Let's give a and b for the weights of r(t,n) and s(n) respectively, given that a+b=1. These weights make the composed performance score as:

$$x = a * r(t,n) + b * s(n)$$

If the aim of the algorithm is to crack a specific password, then more weight would be given to 'a' coefficient, if the aim is to crack as many passwords as possible in a specific time frame, then more weight would be given to 'b' coefficient. Different weight scenarios are considered during performance tests. After having all algorithms graded, a final word list / algorithm would be composed based on the findings, and this list's performance is expected to be superior to base algorithms.

As an example, let's have the following algorithms and corresponding word lists:

- Algorithm / List 1: Combined word lists
- Algorithm / List 2: Combined revealed password lists
- Algorithm / List 3: A custom made password list
- Algorithm / List 4: Lists that use keyboard placement of letters
- Algorithm / List 5: Lists which use frequency analysis

Let's name those algorithms as A1, A2, A3, A4 and A5. p = 1000 is the total number of target passwords, q = 100.000 is maximum try for each of the passwords. After running the script, let's have the corresponding scores for the algorithms:

Algorithm	r(t)	s
A1	62,173,468	17%
A2	36,225,779	37%
A3	28,113,016	16%
A4	58,330,849	64%
A5	36,514,250	32%

After normalizing the scores, we get:

Algorithm	r(t,n)	s(n)
A1	45.21706	26.5625
A2	77.605	57.8125
A3	100	25
A4	48.19579	100
A5	76.9919	50

If we get the a = b = 0.5 and normalize again, we get:

Algorithm	x	x(n)
A1	35.88978	48.43563
A2	67.70875	91.37742
A3	62.5	84.34787
A4	74.0979	100
A5	63.49595	85.69198

And finally, we will construct a final table based on the x(n) scores.

4. Results

4.1. Source and Destination Lists

Both source and destination lists are revealed passwords from real-world. As source list, following files are used. Line number counts are same as password numbers. These lists are used to decide which rules would be applied in the final algorithm.

```

Command Prompt - python
>>> countline_dir(folder = 'C:\\Users\\bt\\OneDrive\\TEZ\\PYTHON\\python_codes\\leaked')
alypaa.txt has 1,384 lines
carders.cc.txt has 1,904 lines
elitehacker.txt has 895 lines
facebook-pastebay.txt has 55 lines
facebook-phished.txt has 2,442 lines
faithwriters.txt has 8,347 lines
hak5.txt has 2,351 lines
hotmail.txt has 8,931 lines
myspace.txt has 37,144 lines
porn-unknown.txt has 8,089 lines
singles.org.txt has 12,234 lines
tuscl.txt has 38,820 lines
>>>
    
```

For the destination (target) lists, the following lists are used:

```

Command Prompt - python
Microsoft Windows [Version 10.0.17763.1339]
(c) 2018 Microsoft Corporation. All rights reserved.

C:\>python
Python 3.8.2 (tags/v3.8.2:7b3ab59, Feb 25 2020, 22:45:29) [MSC v.1916 32 bit (Intel)] on win32
Type "help", "copyright", "credits" or "license()" for more information.
>>> from os import *
>>> countline_dir(folder = 'C:\\Users\\bt\\OneDrive\\TEZ\\PYTHON\\python_codes\\targetlist')
phpbb.txt has 184,389 lines
rockyou.txt has 14,344,391 lines
>>>
    
```

All lists are preprocessed; first checked for uniqueness then randomized. Rockyou.txt list further divided into sub-lists, such that the number of entries in the lists would be same as the number of entries in phpbb.txt (which is 184,389).

```

Command Prompt - python
Microsoft Windows [Version 10.0.17763.1339]
(c) 2018 Microsoft Corporation. All rights reserved.

C:\>python
Python 3.8.2 (tags/v3.8.2:7b3ab59, Feb 25 2020, 22:45:29) [MSC v.1916 32 bit (Intel)] on win32
Type "help", "copyright", "credits" or "license()" for more information.
>>> from os import *
>>> countline_dir(folder = 'C:\\Users\\bt\\OneDrive\\TEZ\\PYTHON\\python_codes\\targetlist2')
rockyou-first 184389 entries.txt has 184,389 lines
rockyou-fourth 184389 entries.txt has 184,389 lines
rockyou-second 184389 entries.txt has 184,389 lines
rockyou-third 184389 entries.txt has 184,389 lines
>>>
    
```

4.2. Source lists unprocessed

In this scenario, source lists are unmodified and they are used as they are. The results are below:

Source: Unprocessed lists, **Target:** phpbb.txt

List name	Match count	Try count	Count per match	Match count		Final Score
				(normalized)	1 / Count per match	
alypaa.txt	583	254,771,047	437,000	8	0.0000022883	70
carders.cc.txt	522	350,446,434	671,353	7	0.0000014895	45
elitehacker.txt	540	164,766,020	305,122	7	0.0000032774	100
facebook-pastebay.txt	11	10,141,083	921,917	0	0.0000010847	33
facebook-phished.txt	267	449,914,131	1,685,072	4	0.0000005934	18
faithwriters.txt	1777	1,529,731,126	860,850	24	0.0000011616	35
hak5.txt	271	433,109,153	1,598,189	4	0.0000006257	19
hotmail.txt	959	1,641,356,684	1,711,529	13	0.0000005843	18
myspace.txt	2545	6,788,598,640	2,667,426	35	0.0000003749	11
porn-unknown.txt	3277	1,475,560,891	450,278	45	0.0000022209	68
singles.org.txt	3048	2,312,201,994	732,350	42	0.0000018655	42
tuscl.txt	7275	6,972,156,311	958,372	100	0.0000010434	32

Source: Unprocessed lists, **Target:** rockyou.txt

List name	Match count	Try count	Count per match	Match count		Final Score
				(normalized)	1 / Count per match	
alypaa.txt	993	19,851,908,939	19,991,852	5	0.0000000500	80
carders.cc.txt	1308	27,310,380,224	20,879,496	7	0.0000000479	77
elitehacker.txt	800	12,837,859,547	15,987,372	4	0.0000000625	100
facebook-pastebay.txt	29	788,940,623	22,204,849	0	0.0000000368	59
facebook-phished.txt	968	35,027,748,990	36,185,691	5	0.0000000276	44
faithwriters.txt	4823	119,710,883,524	24,820,834	24	0.0000000403	64
hak5.txt	568	33,722,905,999	59,371,313	3	0.0000000168	27
hotmail.txt	4042	128,090,376,955	31,889,900	20	0.0000000316	50
myspace.txt	18517	532,435,383,489	28,753,869	94	0.0000000348	56
porn-unknown.txt	6057	116,005,794,772	19,152,352	31	0.0000000522	83
singles.org.txt	7721	175,438,476,189	22,722,248	39	0.0000000440	70
tuscl.txt	19737	556,430,223,902	28,192,239	100	0.0000000355	57
TOTAL		1,757,651,088,159				

In both phpbb and rockyou lists, tuscl list got the maximum score.

4.3. Add character to the items of the list

! ? . @ * 0 and 1 characters are added to the end of the list items. As source list, tuscl.txt is used, which got the maximum score for both target lists, phpbb and rockyou.

The scores of characters:

```

Command Prompt
C:\>python
Python 3.8.2 (tags/v3.8.2:7b3ab59, Feb 25 2020, 22:45:29) [MSC v.1916 32 bit (Intel)] on win32
Type "help", "copyright", "credits" or "license()" for more information.
>>> from os import *
>>> countline_dir(folder = 'C:\\Users\\bt\\OneDrive\\TEZ\\PYTHON\\python_codes\\targetlist2')
rockyou-first 184389 entries.txt has 184,389 lines
rockyou-fourth 184389 entries.txt has 184,389 lines
rockyou-second 184389 entries.txt has 184,389 lines
rockyou-third 184389 entries.txt has 184,389 lines
>>>
    
```

Normalized:

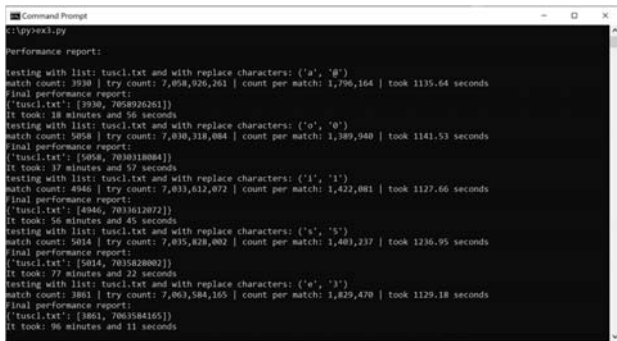
Append Char	Match count	Try count	Count per match	Match count (normalized)	1 / Count per match	1 / Count per match (Normalized)	Final Score
Char !	53	7,156,382,826	135,026,091	4	0.0000000074	4	4
Char ?	2	7,157,904,977	3,578,952,489	0	0.0000000003	0	0
Char :	15	7,157,558,291	477,170,553	1	0.0000000021	1	1
Char @	2	7,157,905,498	3,578,952,749	0	0.0000000003	0	0
Char *	8	7,157,760,386	894,720,948	1	0.0000000011	1	1
Char 0	140	7,154,444,280	51,103,173	10	0.0000000196	10	10
Char 1	1346	7,117,571,384	5,287,943	100	0.0000001891	100	100
TOTAL		50,059,527,642					

Add 1 rule got the maximum score, add 0 rule is following.

4.4. Change character of the items of the list

Change character rules of ('a', '@'), ('o', '0'), ('i', '1'), ('s', '5') and ('e', '3') are applied. As source list, tuscl.txt is used, which got the maximum score for both target lists, phpbbs and rockyou.

The scores of change operations:



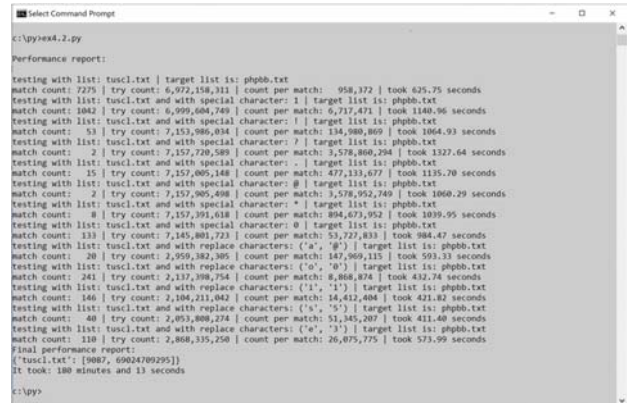
Normalized:

Replace Char	Match count	Try count	Count per match	Match count (normalized)	1 / Count per match	1 / Count per match (Normalized)	Score
('a', '@')	3930	7,058,926,261	1,796,164	78	0.0000005567	77	78
('o', '0')	5058	7,030,318,084	1,389,940	100	0.0000007195	100	100
('i', '1')	4946	7,033,612,072	1,422,081	98	0.0000007032	98	98
('s', '5')	5014	7,035,828,002	1,403,237	99	0.0000007126	99	99
('e', '3')	3861	7,063,584,165	1,829,470	76	0.0000005466	76	76

('o', '0'), ('i', '1') and ('s', '5') change operations got the maximum score.

4.5. All three scenarios combined

When we combine all three scenarios, we get:



Normalized:

OPERATION	match count	try count	count per match	1/gpm	mc (normalized)	1/gpm (normalized)	final score
none	7275	6,972,158,311	958,372	0.000001043485952749	100.00	100.00	100.00
add 1	1042	6,999,604,749	6,717,471	0.00000014886554846527	14.32	14.27	14.29
add !	53	7,153,986,034	134,980,869	0.0000000740845729194	0.73	0.71	0.72
add ?	2	7,157,720,589	3,578,862,295	0.000000027941856281	0.03	0.03	0.03
add :	15	7,157,005,148	477,133,677	0.0000000209584870904	0.21	0.20	0.20
add @	2	7,157,905,498	3,578,952,749	0.00000002794134464	0.03	0.03	0.03
add *	8	7,157,391,618	894,673,952	0.000000011772562226	0.11	0.11	0.11
add 0	133	7,145,801,723	53,727,833	0.0000000186123721512	1.83	1.78	1.81
replace ('a', '@')	20	2,959,382,305	147,969,115	0.000000067581670508	0.27	0.65	0.46
replace ('o', '0')	241	2,137,398,754	8,868,874	0.0000001127538797390	3.31	10.81	7.06
replace ('i', '1')	146	2,104,211,042	14,412,404	0.0000000593846734189	2.01	6.65	4.33
replace ('s', '5')	40	2,053,808,274	51,345,207	0.0000000184701464007	0.55	1.87	1.21
replace ('e', '3')	110	2,868,335,250	26,075,775	0.0000000384977100393	1.51	3.68	2.59
TOTAL	9087	69024709295	7,595,984	0.00000001344850809593	124.91	12.42	68.76

4.6. Selection of Rules, (a, b) = (0.5, 0.5)

If we analyze the marginal utility of operations, we get the following table for values (a, b) = (0.5, 0.5):

OPERATION	match count	try count	count per match	1/gpm	mc (normalized)	1/gpm (normalized)	final score	improvement
none	7275	6,972,158,311	958,372	0.000001043485952749	100.00	100.00	100.00	
add 1	1042	6,999,604,749	6,717,471	0.00000014886554846527	14.32	14.27	14.29	
add !	53	7,153,986,034	134,980,869	0.0000000740845729194	0.73	0.71	0.72	a 0.5
add ?	2	7,157,720,589	3,578,862,295	0.000000027941856281	0.03	0.03	0.03	b 0.5
add :	15	7,157,005,148	477,133,677	0.0000000209584870904	0.21	0.20	0.20	
add @	2	7,157,905,498	3,578,952,749	0.00000002794134464	0.03	0.03	0.03	
add *	8	7,157,391,618	894,673,952	0.000000011772562226	0.11	0.11	0.11	
add 0	133	7,145,801,723	53,727,833	0.0000000186123721512	1.83	1.78	1.81	
replace ('a', '@')	20	2,959,382,305	147,969,115	0.000000067581670508	0.27	0.65	0.46	
replace ('o', '0')	241	2,137,398,754	8,868,874	0.0000001127538797390	3.31	10.81	7.06	
replace ('i', '1')	146	2,104,211,042	14,412,404	0.0000000593846734189	2.01	6.65	4.33	
replace ('s', '5')	40	2,053,808,274	51,345,207	0.0000000184701464007	0.55	1.87	1.21	
replace ('e', '3')	110	2,868,335,250	26,075,775	0.0000000384977100393	1.51	3.68	2.59	
TOTAL	9087	69024709295	7,595,984	0.00000001344850809593	124.91	12.42	68.76	
plus replace ('o', '0')	8558	24153590783	1,633,843	0.0000000465237187868	117.64	57.05	87.79	27.68%
plus replace ('i', '1')	8704	16257801825	1,867,894	0.0000003537372971392	119.64	51.31	85.48	24.31%
plus replace ('s', '5')	8814	19161617079	2,169,970	0.0000004608354617771	121.15	44.17	82.66	20.21%
plus add 0	8947	28271938758	2,996,286	0.000000405345729610	122.88	32.44	77.81	13.16%
plus replace ('s', '5')	8987	28325747072	3,151,838	0.0000003172731831640	123.59	30.41	76.97	11.94%
plus add 1	9040	35479731306	3,924,749	0.0000002547939484464	124.26	24.42	74.34	8.11%
plus replace ('a', '@')	9065	38439115411	4,242,728	0.000000356971142918	124.54	22.59	73.56	6.98%
plus add :	9075	45596120559	5,024,364	0.0000001990309047134	124.74	19.07	71.91	4.58%
plus add *	9083	52753512177	5,807,939	0.0000001721870957353	124.85	16.50	70.68	2.78%
plus add ?	9085	59911232766	6,594,532	0.0000001546410125409	124.88	15.93	69.71	1.37%
plus add @	9087	67069138264	7,380,779	0.0000001348705463057	124.91	12.98	68.95	0.27%

Here maximum score is for the rule 'add 1' (14.29). Second score is for the rule replace ('o' with '0') (7.06), and third score is for the rule replace ('i' with '1') (4.33).

Here, 'add 1' rule increased the overall score by 24.61%, on top of this, replace ('o' with '0') rule increased this number to 27.68%, but replace ('i' with '1') rule decreased this number to 24.31%. So, we can conclude that the rules which must be selected are 'add 1' and replace ('o' with '0').

4.7. Selection of Rules, (a, b) = (0.8, 0.2)

OPERATION	match		count per match	mc		1/rpm	final score	improvement
	count	try count		(normalized)	(normalized)			
none	7275	6,972,158,311	958,372	0.00000104343585952749	100.00	100.00	100.00	
add 1	1042	6,999,604,749	6,717,471	0.0000014886554846527	14.32	14.27	14.31	
add !	53	7,153,986,034	134,980,869	0.0000000740845729184	0.73	0.71	0.72	a
add ?	2	7,157,720,589	3,578,860,295	0.0000000027941856281	0.03	0.03	0.03	b
add *	15	7,157,005,148	477,133,677	0.0000000209504670904	0.21	0.20	0.21	
add @	2	7,157,905,498	3,578,952,749	0.0000000027941134464	0.03	0.03	0.03	
add #	8	7,157,391,618	894,673,952	0.00000001177256226	0.11	0.11	0.11	
add 0	133	7,145,801,723	53,727,833	0.00000001861232723152	1.83	1.78	1.82	
replace ('@', '@')	20	2,959,382,305	147,969,115	0.0000000675816705608	0.27	0.65	0.35	
replace ('@', '0')	241	2,137,398,754	8,868,874	0.000001127538773790	3.31	10.81	4.81	
replace ('@', '1')	146	2,104,211,042	14,412,404	0.000000693846734189	2.01	6.65	2.94	
replace ('@', 'i')	40	2,051,808,274	51,145,207	0.000000194760163407	0.55	1.87	0.81	
replace ('@', 's')	110	2,868,335,250	26,075,775	0.000000383977100393	1.51	3.68	1.94	
TOTAL	9087	69024709295	7,995,984	0.0000013144850808953	124.91	12.62	102.45	
none + add 1	8317	13971763660	1,679,904	0.0000059527204722007	114.32	57.05	102.87	0.41%
plus replace ('@', '0')	8558	14153590783	1,653,843	0.0000060465221378868	117.64	57.95	105.70	3.17%
plus replace ('@', '1')	8704	16257801825	1,867,854	0.000005357372971392	119.64	51.31	105.98	3.44%
plus replace ('@', 'i')	8814	19126137975	2,169,972	0.000004608334617771	121.15	44.17	105.76	3.23%
plus add 0	8947	2627138798	2,936,396	0.0000034055347299610	122.98	32.64	104.91	2.41%
plus replace ('@', 's')	8987	28325747072	3,151,858	0.0000031727318531640	123.59	30.41	104.91	2.40%
plus add !	9040	35479733106	3,924,749	0.0000025479934844644	124.26	24.42	104.29	1.80%
plus replace ('@', '#')	9067	38439115411	4,242,728	0.000002356974142918	124.54	22.59	104.15	1.66%
plus add .	9075	45596120559	5,024,366	0.0000019903090947134	124.74	19.07	103.61	1.13%
plus add !	9083	52753512177	5,807,939	0.0000017217810957353	124.85	16.50	101.18	0.72%
plus add ?	9085	59911232766	6,594,522	0.0000015164101255409	124.88	14.53	101.81	0.35%
plus add @	9087	67069138264	7,380,779	0.0000013548705463057	124.91	12.98	102.52	0.07%

Here maximum score is for the rule 'add 1' (14.31). Second score is for the rule replace ('o' with '0') (4.81), and third score is for the rule replace ('i' with 1) (2.94). Here, 'add 1' rule increased the overall score by 0.41%, on top of this, replace ('o' with '0') rule increased this number to 3.44%, and replace ('i' with 1) rule increased this number to 3.44%. So, we can conclude that the rules which must be selected are 'add 1', replace ('o' with '0') and replace ('i' with '1').

4.8. Selection of Rules, (a, b) = (0.2, 0.8)

OPERATION	match		count per match	mc		1/rpm	final score	improvement
	count	try count		(normalized)	(normalized)			
none	7275	6,972,158,311	958,372	0.00000104343585952749	100.00	100.00	100.00	
add 1	1042	6,999,604,749	6,717,471	0.0000014886554846527	14.32	14.27	14.28	
add !	53	7,153,986,034	134,980,869	0.0000000740845729184	0.73	0.71	0.71	
add ?	2	7,157,720,589	3,578,860,295	0.0000000027941856281	0.03	0.03	0.03	a
add *	15	7,157,005,148	477,133,677	0.0000000209504670904	0.21	0.20	0.20	b
add @	2	7,157,905,498	3,578,952,749	0.0000000027941134464	0.03	0.03	0.03	
add #	8	7,157,391,618	894,673,952	0.00000001177256226	0.11	0.11	0.11	
add 0	133	7,145,801,723	53,727,833	0.00000001861232723152	1.83	1.78	1.79	
replace ('@', '@')	20	2,959,382,305	147,969,115	0.0000000675816705608	0.27	0.65	0.57	
replace ('@', '0')	241	2,137,398,754	8,868,874	0.000001127538773790	3.31	10.81	9.31	
replace ('@', '1')	146	2,104,211,042	14,412,404	0.000000693846734189	2.01	6.65	5.72	
replace ('@', 'i')	40	2,051,808,274	51,145,207	0.000000194760163407	0.55	1.87	1.60	
replace ('@', 's')	110	2,868,335,250	26,075,775	0.000000383977100393	1.51	3.68	3.24	
TOTAL	9087	69024709295	7,995,984	0.0000013144850808953	124.91	12.62	35.07	
none + add 1	8317	13971763660	1,679,904	0.0000059527204722007	114.32	57.05	64.50	95.31%
plus replace ('@', '0')	8558	14153590783	1,653,843	0.0000060465221378868	117.64	57.95	69.89	99.25%
plus replace ('@', '1')	8704	16257801825	1,867,854	0.000005357372971392	119.64	51.31	64.98	85.25%
plus replace ('@', 'i')	8814	19126137975	2,169,972	0.000004608334617771	121.15	44.17	59.56	69.87%
plus add 0	8947	2627138798	2,936,396	0.0000034055347299610	122.98	32.64	50.71	44.57%
plus replace ('@', 's')	8987	28325747072	3,151,858	0.0000031727318531640	123.59	30.41	49.03	39.79%
plus add !	9040	35479733106	3,924,749	0.0000025479934844644	124.26	24.42	44.39	26.55%
plus replace ('@', '#')	9060	38439115411	4,242,728	0.000002356974142918	124.54	22.59	42.98	22.53%
plus add .	9075	45596120559	5,024,366	0.0000019903090947134	124.74	19.07	40.21	14.63%
plus add !	9083	52753512177	5,807,939	0.0000017217810957353	124.85	16.50	38.17	8.83%
plus add ?	9085	59911232766	6,594,522	0.0000015164101255409	124.88	14.53	36.60	4.35%
plus add @	9087	67069138264	7,380,779	0.0000013548705463057	124.91	12.98	35.37	0.84%

Here maximum score is for the rule 'add 1' (14.28). Second score is for the rule replace ('o' with '0') (9.31), and third score is for the rule replace ('i' with 1) (5.72). Here, 'add 1' rule increased the overall score by 95.31%, on top of this, replace ('o' with '0') rule increased this number to 99.25%, but replace ('i' with 1) rule decreased this number to 85.25%. So, we can conclude that the rules which must be selected are 'add 1' and replace ('o' with '0').

5. Conclusion and Future Work

Password usage is a very prevalent security control used to protect information assets. To assess the effectiveness of this control, security professionals may use password cracking tools and algorithms. This study proposes a method to assess success criterion of different password cracking algorithms; and based on the assessment scores, proposes another method to combine these algorithms in a specific order, to increase the likelihood of password cracking in a most efficient and effective manner. In this study it is shown that, based on the aim of the password cracking (to crack a specific password vs to crack as many passwords as possible in a specific time frame), different strategies may be followed. And this study offers a way to determine which algorithms / lists to include in the final cracking endeavor, and which ones to exclude.

This work studied add and replace rules. Further study may include keyboard patterns, frequency analysis, etc. Also this study used (a,b) weights of (0.2, 0.8), (0.5, 0.5) and (0.8, 0.2). Different weights, different source lists, different target lists, different add / replace rules, different password try cutoff points may be tried. If data can be obtained, user demographics, target system features, password length, etc. may be used as a decision point in the final algorithm.

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